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Non-decreasing Sequences

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Non-Decreasing Sequences

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Abstract

Non-decreasing sequences are a generalization of binary covering arrays, which has made research on non-decreasing sequences important in both math and computer science. The goal of this research is to find properties of these non-decreasing sequences as the variables d , s , and t change. The goal is also to explore methods for creating a maximum length non-decreasing sequence for a given strength and size set. Through our research, we discovered and proved basic properties of these non-decreasing sequences. In addition to this, we can describe a method we used while trying to find the maximum length of a sequence.

Definitions and Notation

- Let S be a set of s elements
- The **strength** of non-decreasing sequence is the amount of subsets whose union we consider, and is represented using d
- A **non-decreasing sequence** of strength d is a sequence of non-empty subsets, $\{S_1, S_2, \dots, S_t\}$, where the union of any d previous subsets does not contain any subsequent subset
- The number of subsets in a non-decreasing sequence is called the **length**, t
- $NDS(d, s, t)$ is the set of non-decreasing sequences with strength d , s elements and length t
- $NDST(d, s)$ is the maximum t such that $NDS(d, s, t)$ is non-empty
- Let r_j be the number of elements in the subset S_j

Binary Arrays

- Represent a non-decreasing sequence using an $s \times t$ binary array
- Rows represent elements of S
- Columns represent subsets of non-decreasing sequence

| | |
|-----|-------|
| | S_i |
| 1 | 0 |
| ⋮ | ⋮ |
| k | 1 |
| ⋮ | ⋮ |
| s | 0 |

| | | | | | |
|---|-------|-------|-------|-------|-------|
| | S_1 | S_2 | S_3 | S_4 | S_5 |
| 1 | 1 | 0 | 0 | 0 | 1 |
| 2 | 0 | 1 | 0 | 0 | 1 |
| 3 | 0 | 0 | 1 | 0 | 1 |
| 4 | 0 | 0 | 0 | 1 | 0 |

Basic Results

Theorem 1-Permuting rows in a binary array gives another $NDS(d, s, t)$.

Theorem 2-If the union of any d subsets contain all elements in S , no subsets can be added to the sequence.

Theorem 3-Every subset in $NDS(d, s, t)$ must be distinct for $d \geq 1$.

Theorem 4- $NDS(d, s, t) \subseteq NDS(d, s + 1, t)$

Corollary 5- $NDST(d, ks) \geq kNDST(d, s)$, where $k \in \mathbb{Z}$.

| | | | | | |
|---|-------|-------|-------|-------|-------|
| | S_1 | S_2 | S_3 | S_4 | S_5 |
| 1 | 1 | 0 | 0 | 0 | 1 |
| 2 | 0 | 1 | 0 | 0 | 1 |
| 3 | 0 | 0 | 1 | 0 | 1 |
| 4 | 0 | 0 | 0 | 1 | 0 |

A

| | | |
|---|---|---|
| A | 0 | 0 |
| 0 | A | 0 |
| 0 | 0 | A |

Block array

Standard Sequence

Theorem 6-There exists an $NDS(d, s, t)$ where the first s subsets are of size 1. We call this a **standard non-decreasing sequence**.

| | | | | | | |
|---|---|---|---|---|---|-----|
| | 1 | 2 | 3 | 4 | 5 | 6 |
| 1 | 1 | 0 | 0 | 0 | 0 | ... |
| 2 | 0 | 1 | 0 | 0 | 0 | ... |
| 3 | 0 | 0 | 1 | 0 | 0 | ... |
| 4 | 0 | 0 | 0 | 1 | 0 | ... |
| 5 | 0 | 0 | 0 | 0 | 1 | ... |
| ⋮ | ⋮ | ⋮ | ⋮ | ⋮ | ⋮ | ⋮ |

Theorem 7-A standard non-decreasing sequence of strength d does not have any subsets of size $1 < r \leq d$.

Theorem 8- In a standard non-decreasing sequence, any subset S_j of size $r_j = d + 1$, may contain at most 1 element from any previous subset S_i .

Theorem 9-In a standard non-decreasing sequence, any subset S_j of size $r_j \geq d + 1$ must contain at least d elements that differ from any previous subset S_i .

Bounds

- Gives range for $NDST(d, s)$
- Lower bound is the length of sequence constructed for a given d, s
- Upper bound initially $2^s - 1$, number of nonempty subsets possible for any set S with s elements
- Upper bound decreased using Theorems 7, 8, and 9

| s | Lower Bound | Found Upper Bound | $2^s - 1$ |
|-----|-------------|-------------------|-----------|
| 1 | 1 | 1 | 1 |
| 2 | 2 | 2 | 3 |
| 3 | 4 | 4 | 7 |
| 4 | 5 | 5 | 15 |
| 5 | 7 | 7 | 31 |
| 6 | 11 | 13 | 63 |
| 7 | 15 | 20 | 127 |

Table 1: Bounds for $d = 2$

Future Work

- Find exact formula for $NDST(d, s)$
- Find different computational methods
- Find relation to binary covering arrays
- Effect of permuting columns
- Find bounds for larger d and s values

References

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